

Injectives Hulls and Projective Covers in Categories of Generalized Uniform Hypergraphs

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Author's contribution

The sole author designed, analyzed, interpreted and prepared the manuscript.

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Abstract

We construct injective hulls and projective covers in categories of generalized uniform hypergraphs which generalizes the constructions in the category of quivers and the category of undirected graphs. While the constructions are not functorial, they are "sub-functorial", meaning they are subobjects of functorial injective and projective refinements.

Keywords: Graph; hypergraph; uniform hypergraph; injective hull; projective cover.

1 Introduction

In Grilliette $[1]$, the construction of injective hulls and projective covers of directed graphs² are given. We provide a new approach to this construction by considering categories of generalized *k*-uniform hypergraphs. This construction also gives us injective hulls and projective covers of

^{}Corresponding author: E-mail: martin.schmidt@wnc.edu* ² also calle[d](#page-11-0) quivers

undirected graphs, uniform hypergraphs, as well as a multitude of other kinds of "graphs" which have a set of vertices, edges, and incidence maps.

We introduce categories of (*X, M*)-graphs where *M* is a monoid and *X* is a right *M*-set. An (*X, M*) graph *G* consists of a set of arcs $G(A)$ and a set of vertices $G(V)$ such that each arc is incident to #*X*-vertices (multiplicities allowed) and *M* informs the type of cohesivity between the vertices. The categories of (X, M) -graphs introduced in this paper are able to describe the types of incidence in the various definitions of graphs and hypergraphs by taking the monoid *M* to be a submonoid of endomaps on a set *X*. Thus, when *X* is a two-element set, the categories of (*X, M*)-graphs generalize the various categories of graphs and undirected graphs found in Bumby and Latch [2].

By separating syntax $((X, M)$ -graph theories) from semantics $((X, M)$ -graph categories), functorial constructions between (X, M) -graph categories are induced from the appropriate morphisms between theories. In particular, the constructions of injective hulls and projective covers (Section 3) can be obtained by using the obvious morphisms of monoid actions as well as obvious interpretatio[ns](#page-11-1) of (X, M) -graph theories.

2 Categories of (*X, M*)**-Graphs**

We begin with a definition.

Definition 2.1. Let *M* be a monoid and *X* a right *M*-set. The theory for (X, M) -graphs, $\mathbb{G}_{(X, M)}$, has two objects *V* and *A* with homsets given by

$$
\mathbb{G}_{(X,M)}(V,A) := X,
$$

\n
$$
\mathbb{G}_{(X,M)}(A,V) := \emptyset,
$$

\n
$$
\mathbb{G}_{(X,M)}(V,V) := \{\mathrm{id}_V\},
$$

\n
$$
\mathbb{G}_{(X,M)}(A,A) := M.
$$

Composition is defined as $m \circ x = x.m$ (the right-action via *M*), $m \circ m' = m'm$ (monoid operation of *M*). The category of (X, M) -graphs is defined to be the category of presheaves $\mathbb{G}_{(X,M)} :=$ $[\mathbb{G}_{(X,M)}^{op}, \mathbf{Set}]$.

We represent the theory for (*X, M*)-graphs and and reflexive *M*-graphs as follows.

$$
V \mathop{\longrightarrow} A \bigcap M
$$

By definition, an (X, M) -graph $G: \mathbb{G}_{(X,M)}^{op} \to \mathbf{Set}$ has a set of vertices $G(V)$ and a set of arcs $G(A)$ along with right-actions for each morphism in $\mathbb{G}_{(X,M)}$. For example, $x: V \to A$ in $\mathbb{G}_{(X,M)}$ yields a set map $G(x)$: $G(A) \rightarrow G(V)$ which takes an arc $\alpha \in G(A)$ to $\alpha.x := G(x)(\alpha)$ which we think of as its *x*-incidence.³ For an element *m* in the monoid *M*, the corresponding morphism $m: A \rightarrow A$ in $G(X,M)$ yields a right-action $\alpha.m := G(m)(\alpha)$ which we think of as the *m*-associated partner of α .

Each (X, M) -graph *G* induces a set map $\partial_G : G(A) \to G(V)^X$ such that $\partial_G(\alpha) : X \to G(V)$ is the parametriz[ed](#page-1-0) incidence of α , i.e., $\alpha x = \partial_G(\alpha)(x)$. The *x*-incidence can be recovered from a parametrized incidence by precomposition of the map α : $1 \rightarrow X$ which names the element *x* in *X*. Observe that the *m*-associated partner of an arc α in G has the parametrized incidence such that

³Note that we use the categorical notation of evaluation of a presheaf as a functor for the set of vertices $G(V)$ and set of arcs $G(A)$ rather than the conventional graph theoretic $V(G)$ and $E(G)$ for the vertex set and edge set.

the following commutes

$$
X \xrightarrow{\langle \text{id}_X, \lceil m \rceil \rangle} X \times M \xrightarrow{\text{action}} X \xrightarrow{\partial_G(\alpha)} G(V).
$$

Example 2.1. *Let X be a set and M a submonoid of endomaps* End(*X*)*. The right-action of* $M \subseteq \text{End}(X)$ *on X is given by evaluation, e.g.* $x.f := f(x)$ *. When X is a two-element set the categories* $\widehat{\mathbb{G}}_{(\text{id}X)}$ *,X*) *and* $\widehat{\mathbb{G}}_{(\text{Aut}(X),X)}$ *are the categories of quivers [1] and symmetric directed graphs [2]. For example:*

- *1. When* $X = \emptyset$ *and M is the trivial monoid, the category* $\widehat{\mathbb{G}}_{(X,M)}$ *is* **Set** \times **Set***.*
- 2. When $X = 1$ is a [on](#page-11-0)e element set and M is the trivial monoid, $\widehat{\mathbb{G}}_{(X,M)}$ is the category of *bouquets i.e., the category of presheaves on* $V \stackrel{s}{\rightarrow} A$ ([3], p 18).
- *3. When X* = *{s, t} and M is the submonoid of endomorphisms which exchanges s and t, the categories* $\widehat{\mathbb{G}}_{(X,M)}$ *is the category of undirected graphs with involution described in Brown et al. [4].*

T[he](#page-12-0) following is an example of such an (X, M) *-graph where* $i: X \rightarrow X$ *denotes the non-trivial automap.*

To connect this definition to undirected graphs, we identify edges which are i-pairs and define the set of edges G(*E*) *as the quotient of the set of arrows G*(*A*) *under this automorphism defined by the i-action.*⁴ *There is an incidence operator* $\partial : G(E) \to G(V)^2$ *which defines for an i-pair the set of boundaries. Then an undirected representation for G can be given as*

We have placed a 2 *in the loop which came from the 2-loop* $\beta_0 \sim \beta_1$ *even though the quotient has identified them.*

2.1 Nerve-Realization Adjunctions

The symbols and notation in this section follow from Applegate and Tierney [5] and Riehl [6].

Let $I: \mathbb{T} \to M$ be functor from a small category \mathbb{T} to a cocomplete category M. Since the Yoneda

⁴In the subsequent, we reserve the term edge for the equivalence class of arcs under the group $s(X)$.

embedding $y: T \to \hat{\mathbb{T}}$ is the free cocompletion of a small category there is a essentially unique adjunction $R \dashv N$: $M \to \hat{T}$, called the nerve-realization adjunction, such that $Ry \cong I$.

The nerve and realization functors are given on objects by $N(m) = M(I(-), m)$, $R(X) = \text{colim}_{(c,\varphi) \in \int F} I(c)$ respectively, where $\int F$ is the category of elements of *X* ([5], Section 2, pp $124-126$).⁵

We call a functor $I: \mathbb{T} \to M$ from a small category to a cocomplete category an interpretation functor. The category \mathbb{T} is called the theory for *I* and *M* the modeling category for *I*. An interpretati[on](#page-3-0) $I: \mathbb{T} \to M$ is dense, i.e., for each *M*-object *m* is isomorphic to the [co](#page-12-2)limit of the diagram $I \downarrow m \rightarrow M$, $(c, \varphi) \mapsto I(c)$, if and only if the nerve $N: M \rightarrow \hat{\mathbb{T}}$ is full and faithful ([7], Section X.6, p 245). When the right adjoint (resp. left adjoint) is full and faithful we call the adjunction reflective (resp. coreflective).⁶

A functor $F: C \to C'$ between small categories C and C' induces an essential geometric morphism $F_1 + F^* + F_* : \widehat{C} \to \widehat{C'}$ $F_1 + F^* + F_* : \widehat{C} \to \widehat{C'}$ $F_1 + F^* + F_* : \widehat{C} \to \widehat{C'}$ where F_1 is the realization of $y_{C'} \circ F$, F^* is the nerve of $y_{C'} \circ F$ and F_* is the nerve of $F^* \circ y_{C'}$ where $y_C: C \to \widehat{C}$ and $y_{C'}: C' \to \widehat{C'}$ are the Yoneda embeddings (see Reyes a[n](#page-3-1)d Reyes $[3]$ pp 194-198, Section 4). On objects *W* in *C* and *Z* in C' , the functors are given by

$$
F_!(W) := \text{colim}_{(C,c) \in \int W} y_{C'} F(C),
$$

\n
$$
F^*(Z) := \widehat{C'}(y_{C'} F(-), Z),
$$

\n
$$
F_*(W) := \widehat{C}(F^* y_{C'}(-), W)
$$

where ∫ *A* is the category of elements for *A* ([6], Section 2.4). In the subsequent, we denote the representables for the vertex and arc object by *V* and *A* respectively.

Consider the (X, M) -graph theory $\mathbb{G}_{(\varnothing, 1)}$, i.e., the discrete category with two objects *V* and *A*. Then for an (X, M) -graph theory $\mathbb{G}_{(X,M)}$, there is the inclusion functor $\iota: \mathbb{G}_{(\emptyset,1)} \to \mathbb{G}_{(X,M)}$. Thus there is an essential geometric morphism $\iota_! \dashv \iota^* \dashv \iota_* : \mathbf{Set}^2 \to \widehat{\mathbb{G}}_{(X,M)}$. The *[ι](#page-12-4)*-extension $\iota_!$ takes the pair of sets $(S(V), S(A))$ to the coproduct $\bigsqcup_{S(V)} \underline{V} \sqcup \bigsqcup_{S(A)} \underline{A}$ since the category of elements for $(S(V), S(A))$ lacks internal cohesion. The *ι*-restriction ι^* takes an (X, M) -graph *G* to the pair of sets $(G(V), G(A))$. By Riehl [6] (Proposition 3.3.9), it creates all limits and colimits. The *ι*-coextension ι_* sends $(S(V), S(A))$ to the (X, M) -graph with vertex set $S(V)$ and arc set $\textbf{Set}^2((X, |M|), (S(V), S(A))) = S(V)^X \times S(A)^{|M|}$ where |M| is the underlying set of M. The rightactions are given by $(f, s) \cdot x = f(x)$, $(f, s) \cdot m = (f \circ m, s \cdot m)$ where $m \colon X \to X$ is the right-action map by $m \in M$ and $s.m: |M| \rightarrow S(A)$ $s.m: |M| \rightarrow S(A)$ is defined $s.m(m') := s(mm')$.

The counit $\varepsilon: \iota_! \iota^* \Rightarrow \text{id}$ of the adjunction $\iota_! \dashv \iota^*$ on a component $\varepsilon_G: \bigsqcup_{G(V)} \underline{V} \sqcup \bigsqcup_{G(A)} \underline{A} \to G$ is the epimorphism induced by the classification maps $v: \underline{V} \to G$ and $\alpha: \underline{A} \to G$ for vertices $v \in G(V)$ and arcs $\alpha \in G(A)$. Therefore, the *ι*-restriction functor is faithful and thus $\iota^* \colon \widehat{\mathbb{G}}_{(X,M)} \to \mathbf{Set}^2$ is monadic ([8], p. 227)

The unit $\eta: id \Rightarrow \iota_* \iota^*$ of the adjunction $\iota^* \dashv \iota_*$ on a component $\eta_G: G \to \iota_* \iota^*(G)$ is the identity on vertices and sends arc $\alpha \in G(A)$ to $(\partial_G(\alpha), \overline{\alpha})$ where $\overline{\alpha}$: $|M| \to G(A)$ is the constant map. Thus for each (X, M) -graph *G* the component η_G is a monomorphism.

⁵In [5], the nerve functor is called the singular functor.

⁶since it implies *M* is equivalent to a reflective (resp. coreflective) subcategory of $\hat{\mathbb{T}}$

3 Injective and Projective (*X, M*)**-Graphs**

We set Proj := $\iota_i \iota^*$ and Inj := $\iota_* \iota^*$ where $\iota: \mathbb{G}_{(\varnothing,X)} \to \mathbb{G}_{(X,M)}$ is the functor given above. Then since adjunctions are closed under composition, we have

$$
\text{Proj } \exists \text{ Inj}: \widehat{\mathbb{G}}_{(X,M)} \to \widehat{\mathbb{G}}_{(X,M)}.
$$

We will show in this section that the natural transformations ε : Proj \Rightarrow id and η : id \Rightarrow Inj can be thought of as the functorial projective and injective refinements for non-initial (*X, M*)-graphs.

We first characterize the class of injective and projective objects. Recall that an object *Q* in a category *C* is injective provided for each monomorphism $m: A \rightarrow B$ and morphism $f: A \rightarrow Q$ there exists a morphism (not necessarily unique) $k: B \to Q$ such that $f = km$. Dually, an object P in *C* is (regular) projective⁷ provided for each (regular) epimorphism $e: B \to A$ and each morphism $f: P \to A$ there is a morphism $k: P \to B$ such that $f = ek$.⁸

Proposition 3.1. *A* (*X, M*)*-graph Q is injective if and only if Q is non-initial and for each set* map $f: X \to Q(V)$, the[re](#page-4-0) is an arc $\alpha \in Q(A)$ such that the incidence map $\partial_Q(\alpha)$ is equal to f.

Proof. Suppose *Q* is injective and consider the set map $f: X \to Q(V)$ $f: X \to Q(V)$ $f: X \to Q(V)$. This is equivalent to giving an (X, M) -graph morphism \overline{f} : $\bigsqcup_{x \in X} \underline{V} \to Q$. Consider the inclusion m : $\bigsqcup_{x \in X} \underline{V} \to \underline{A}$ induced by the morphisms \underline{x} : $\underline{V} \rightarrow \underline{A}$. Since *Q* is injective, there is a morphism α : $\underline{A} \rightarrow I$ such that $\alpha m = \overline{f}$. By Yoneda, this is equivalent to an arc $\alpha \in I(A)$ with incidence map $\partial_I(\alpha) = f$.

Conversely, let $f: G \hookrightarrow H$ be a monomorphism and $g: G \to Q$ a morphism of (X, M) -graphs. Since *Q* is non-initial, there is a vertex $v \in Q(V)$. Each arc α in *H* has incidence $\partial_H(\alpha): X \to H(V) \cong$ $f_V(G(V)) \sqcup H(V) \backslash f_V(G(V))$ where $f_V(G(V))$ is the image of the vertices in G under f. For each arc *α* in *H* not in the image of f_A , let $j_\alpha: X \to Q(V)$ be the set map $[g_V, !] \circ \partial_G(\alpha)$ given by universal property of the disjoint union

$$
f_V(G(V)) \cong f_V^{-1}(G(V)) \xrightarrow{g_V} g_V(G(V))
$$

$$
\downarrow \qquad \qquad \downarrow
$$

$$
X \xrightarrow{\partial_G(\alpha)} f_V(G(V)) \sqcup H(V) \setminus f_V(G(V)) \xrightarrow{[g_V, !]} \qquad \qquad Q(V)
$$

$$
\uparrow \qquad \qquad \downarrow
$$

$$
H(V) \setminus f_V(G(V)) \xrightarrow{!} \qquad \qquad \downarrow
$$

$$
f_V
$$

Thus by assumption, we may choose an arc $[\alpha] \in Q(A)$ with incidence equal to j_{α} . We define the following maps $h_V: H(V) \to Q(V)$ and $h_A: H(A) \to Q(A)$

$$
h_V(w) := \begin{cases} g_V(u) & \text{if } \exists u \in G(V), f_V(u) = w \\ v & \text{if } \forall u \in G(V), f_V(u) \neq w \end{cases}
$$

$$
h_A(\alpha) := \begin{cases} g_A(\beta) & \text{if } \exists \beta \in G(A), f_A(\beta) = \alpha \\ [\alpha] & \text{if } \forall \beta \in G(A), f_A(\beta) \neq \alpha \end{cases}
$$

By construction this defines a morphism $h: H \to Q$ such that $h \circ f = g$. Therefore Q is injective. П

⁷Note that since regular epimorphisms are equivalent to epimorphisms in categories of presheaves, a regular projective object is equivalent to a projective object.

⁸The results in this section generalize the results of Grilliette $[1]$ and Williams $[9]$.

Corollary 3.1. *The class of injective objects in* $\mathbb{G}_{(X,M)}$ *is precisely the class of non-initial split subobjects*⁹ *of objects in the essential image of the functor* Inj*.*

Proof. Let *Q* be an injective object in $\mathbb{G}(X,M)$. Hence *Q* is non-initial and thus by the previous lemma, we have Inj(*Q*) is an injective object. Then since $\eta_Q: Q \to \text{Inj}(Q)$ is a monomorphism, there mu[st](#page-5-0) be a split epimorphism $r: \text{Inj}(Q) \to Q$ such that $r\eta_Q = id$ by the property of *Q* being injective.

We also have the dual argument that the class of projective objects in the category of (X, M) -graphs is precisely the split quotients of objects in the essential image of Proj.

Proposition 3.2. *A* (*X, M*)*-graph P is projective if and only if it is a coproduct of representables.*

Proof. Suppose *P* is projective. Since ε_P : Proj $(P) \rightarrow P$ is an epimorphism, there must exist a section $s: P \to \text{Proj}(P)$ such that $\varepsilon_P s = \text{id}$ by the property that P is projective. Since $\text{Proj}(P)$ is a coproduct of representables and *s* is a split monomorphism, *P* must also be a coproduct of representables. Conversely, representables \underline{V} and \underline{A} in a category of presheaves are always projective. Since projective objects are closed under coproducts, the reverse condition is also true.

Corollary 3.2. *The class of projective objects in* $\widehat{\mathbb{G}}_{(X,M)}$ *is the object class of the essential image of the functor* Proj*.*

Proof. Given a coproduct of representables $\bigsqcup_{S} \underline{V} \sqcup \bigsqcup_{T} \underline{A}$, let *H* be the (X, M) -graph with vertex set $H(V) := S$ and arc set $H(A) = T$. Take some $s \in S$ and define right-actions $t.x := s$ and $t.m = t$ for each $t \in T$, $x \in X$ and $m \in M$. Then $\text{Proj}(H) \cong \bigsqcup_{S} \underline{V} \sqcup \bigsqcup_{T} \underline{A}$. \Box

Next, we construct injective hulls and projective covers for (X, M) -graphs. Recall that a monomorphism $i: G \to \tilde{G}$ is essential provided for each morphism $h: G \to H$ such that hi is a monomorphism implies *h* is a monomorphism. An injective hull of an object *G* is an essential monomorphism $i: G \to \widetilde{G}$ where \tilde{G} is injective. Dually, an epimorphism $e: \overline{G} \to G$ is essential provided for each morphism $h: H \to G$ such that *eh* is an epimorphism implies *h* is an epimorphism. A projective cover of an object *G* is an essential epimorphism $e: \overline{G} \to G$ where \overline{G} is projective.

In the case of the initial (X, M) -graph 0, it is straightforward to verify the terminal morphism $0 \to 1$ is the injective hull. For a non-initial (X, M) -graph *G* we define \widetilde{G} to be the (X, M) -graph with vertex set $\widetilde{G}(V) := G(V)$ and arcs set $\widetilde{G}(A) := G(A) \sqcup \{f : X \to G(V) | \forall \alpha \in G(A), \partial_G(\alpha) \neq f\}$ with the obvious right-action. Then \widetilde{G} is an injective object and there is an obvious inclusion $i: G \rightarrow$ \widetilde{G} . To show that it is essential, let $h: \widetilde{G} \to H$ be a morphism such that hi is a monomorphism. Then since *i* is bijective on vertices, h_V must be injective. On arcs, it is enough to show that h_A is injective on $G(A)\backslash G(A)$. However, this is trivial since there is only one arc $f \in G(A)\backslash G(A)$ with incidence $f: X \to G(V)$. Hence *h* is a monomorphism.

For the projective cover, we define $\overline{G} := \bigsqcup_S \underline{V} \sqcup \bigsqcup_T \underline{A}$ where

 $S := \{v \in G(V) | \forall \alpha \in G(A), \forall x \in X, \alpha \cdot x \neq v\}$, i.e., S is the set of isolated vertices in G, and T is a generating subset of $G(A)$ for the right *M*-action $G(A) \times M \rightarrow G(A)$ of minimal cardinality, i.e., for each $\alpha \in G(A)$ there exists a $\beta \in T$ and an element $m \in M$ such that $\beta.m = \alpha$. Since T generates $G(A)$ under the right-action of M and S is the set of vertices of G which are not incident to an arc, the restriction of $\varepsilon_{G|\overline{G}}$: $\overline{G} \to G$ is an epimorphism. It is clear that if $h: H \to \overline{G}$ is a

 \Box

 \Box

⁹An object *H* is a split subobject of *G* provided it admits a split monomorphism $s: H \to G$.

morphism such that *he* is an epimorphism, then *h* must be an epimorphism since \overline{G} is a coproduct of representables of minimal size.

Note that these assignments \tilde{G} and \overline{G} do not extend to functors since there is choice involved. However, by construction we see that both the injective hull \tilde{G} and projective cover \overline{G} embed into $Inj(G)$ and $Proj(G)$ which are functorial constructions.

4 Applications

We follow the definition given in Jakel [10].

Definition 4.1. Let $F:$ **Set** \rightarrow **Set** be an endofunctor. The category of F -graphs G_F is defined to be the comma category $G_F := \mathbf{Set} \downarrow F$.

In other words, an *F*-graph $G = (G(E), G(V), \partial_G)$ $G = (G(E), G(V), \partial_G)$ $G = (G(E), G(V), \partial_G)$ consists of a set of edges $G(E)$, a set of vertices $G(V)$ and an incidence map $\partial_G: G(E) \to F(G(V))$. A morphism $(f_E, f_V): (G(E), G(V), \partial_G) \to$ $(H(E), H(V), \partial_H)$ is a pair of set maps $f_E : G(E) \to H(E)$ and $f_V : G(V) \to H(V)$ such that the following square commutes

$$
G(E) \xrightarrow{f_E} H(E)
$$

$$
\partial_G \downarrow \qquad \qquad \downarrow \partial_H
$$

$$
F(G(V)) \xrightarrow{F(f_V)} F(H(V)).
$$

It is well-known that the category of *F*-graphs is cocomplete with the forgetful functor $U: G_F \rightarrow$ $\textbf{Set} \times \textbf{Set}$ creating colimits [10].

Let $\mathbb{G}_{(X,M)}$ be a theory for (X,M) -graphs and *q* an element in $F(X)$ such that $F(m)(q) = q$ for each $m \in M$ where $m: X \to X$ is the right-action map. We define $I(V) := (\emptyset, 1, 1)$ *,* and $I(A) := (1, X, \lceil q \rceil)$ where $l_1 : \emptyset \to 1$ is the initial map and $\lceil x \rceil : 1 \to X$ the set map with evaluation at $x \in X$. On morphisms, w[e se](#page-12-5)t

$$
(x: V \to A) \quad \mapsto \quad I(x) := (1_1, x^*) \colon (\emptyset, 1, 1_1) \to (1, X, \lceil q \rceil),
$$

$$
(m: A \to A) \quad \mapsto \quad I(m) := (\text{id}_1, F(m)) \colon (1, X, \lceil q \rceil) \to (1, X, \lceil q \rceil).
$$

Verification that $I: \mathbb{G}_{(X,M)} \to G_F$ is a well-defined interpretation functor is straightforward.

4.1 The Category of Hypergraphs

We recall that a hypergraph $H = (H(V), H(E), \varphi)$ consists of a set of vertices $H(V)$, a set of edges $H(E)$ and an incidence map $\varphi: H(E) \to P(H(V))$ where $P:$ **Set** \to **Set** is the covariant power-set functor. In other words, we allow infinite vertex and edge sets, multiple edges, loops, empty edges and empty vertices.¹⁰ In other words the category of hypergraphs H is the category of P -graphs.

Let *X* be a set and apply the definition for the interpretation given above in 4 for $\mathbb{G}_{(X,\text{Aut}(X))}$ with

¹⁰An empty vertex is a vertex not incident to any edge in $H(E)$. An empty edge is an edge e such that $\varphi(e) = \varnothing$ [.](#page-6-1)

q := *X* in *P*(*X*). Note that for each automap $\sigma: X \to X$, *P*(σ) is the identity map. Thus the interpretation $I: \mathbb{G}_{(X,\text{Aut}(X))} \to H$ defined in Section 4 is a well-defined functor.

The nerve $N: H \to \widehat{\mathbb{G}}_{(X,\text{Aut}(X))}$ induced by *I* takes a hypergraph $H = (H(E), H(V), \varphi)$ to the $(X, \text{Aut}(X))$ -graph $N(H)$ with vertex and arc set given by

$$
N(H)(V) = H(I(V), H) = H(V),
$$

\n
$$
N(H)(A) = H(I(A), H) = \left\{ (\beta, f) \in H(E) \times H(V)^{X} \middle| P(f) = \varphi(\beta) \right\}
$$

Notice that in the case a hyperedge e has less than $\#X$ incidence vertices the nerve creates multiple edges and if a hyperedge has more than $\#X$ incidence vertices there is no arc in the correponding $(X, \text{Aut}(X))$ -graph given by the nerve (see Example 4.1 below).

The realization $R: \widehat{\mathbb{G}}_{(X,\text{Aut}(X))} \to H$ sends a $(X,\text{Aut}(X))$ -graph *G* to the hypergraph $R(G)$ = $(R(G)(E), R(G)(V), \psi)$ with vertex, edge sets and incidence map given by

$$
R(G)(V) = G(V),
$$

\n
$$
R(G)(E) = G(A)/\sim, \quad (\sim \text{ induced by Aut}(X)),
$$

\n
$$
\psi \colon R(G)(E) \to P(R(G)(V)), \quad |\gamma| \mapsto \{ v \in G(V) \mid \exists x \in X, \gamma.x = v \}
$$

For a $(X, \text{Aut}(X))$ -graph morphism $f: G \to G'$, the hypergraph morphism $R(f): R(G) \to R(G')$ has $R(f)_V := f_V$ and $R(f)_E := [f_A]$ where $[f_A]: \frac{G(A)}{\sim} \to \frac{G'(A)}{\sim}$ is induced by the quotient.

Example 4.1.

- 1. Let $X = \{a, b, c\}$ and consider the hypergraph *H* with two vertices 0 and 1 and one hyperedge *α between them. Then the nerve* $N(H)$ *has two vertices* 0 *and* 1 *and arc set* $N(H)(A) =$ *{*001*,* 010*,* 100*,* 011*,* 101*,* 110*} with a pair of* Aut(*X*)*-partners* 001 *∼* 010 *∼* 100 *and* 011 *∼* 101 *∼* 110*. The realization identifies the* Aut(*X*)*-partners and thus RN*(*H*) *has two vertices* 0 *and* 1 *and two edges* [001] *and* [011] *between them. The counit* $\varepsilon_H : RN(H) \to H$ *is bijective on vertex set and sends* [001] *and* [011] *to* α *.*
- *2. Let* $X = \{a, b, c\}$ *and consider the hypergraph H with four vertices and one hyperedge* α *connecting them. The nerve* $N(H)$ *has vertex set equal to* $H(V)$ *but empty arc set* $N(H)(A)$ = \emptyset *. The counit* $\varepsilon_H : RN(H) \to H$ *is the inclusion of vertices.*

We are able to use the adjunction above to classify the projective objects in the category of hypergraphs. Recall that a right adjoint functor is faithful if and only if the counit is an epimorphism.

Lemma 4.2. Let *X* be a set with cardinality κ greater than 1. Then the nerve $N: H \to \widehat{\mathbb{G}}_{(X,\text{Aut}(X))}$ *of the interpretation* $I: \mathbb{G}_{(X,\text{Aut}(X))} \to H$ *is faithful on the full subcategory* H_{κ} *consisting of hypergraphs H* with $\max_{e \in H(E)} \varphi(e)$ *at most of cardinality* κ *.*

Proof. Given hypergraph *H* in the subcategory H_k the counit $\varepsilon_H : RN(H) \to H$ of the adjunction $R \dashv N$: $H \to \widehat{\mathbb{G}}_{(X,\text{Aut}(X))}$ at component *H* is an epimorphism since for each hyperedge $e \in H(E)$, there is a $(X, \text{Aut}(X))$ -graph morphism $f: I(A) \to H$ such that f_E takes the lone hyperedge in *I*(*A*) to *e*. \Box

Let $I(V)$ be the hypergraph with one vertex and no hyperedges. For each cardinal number k , let $E_k := I(A)$ be the hypergraph where $I: \mathcal{S} \mathbb{G}_k \to H$ is the interpretation functor above.

Lemma 4.3. *The proper class of objects consisting of the vertex object* $I(V)$ *and hyperedge objects* $(E_k)_{k \in \mathbf{Set}}$ *is a family of separators for the category of hypergraphs.*

Proof. Let $f, g: H \to H'$ be distinct hypergraph morphisms. Let κ be the maximum of the cardinalities such that *H* and *H'* are in the subcategory H_{κ} . By the lemma above, the nerve *N*: $H \to \widehat{\mathbb{G}}_{(X, \text{Aut}(X))}$ is faithful on H_{κ} . Thus we have $N(f) \neq N(g)$. Therefore, either $I(V)$ separates *f* and *g* or $I(A) = E_k$ separates *f* and *g* by definition of the nerve functor. separates *f* and *g* or $I(A) = E_k$ separates *f* and *g* by definition of the nerve functor.

Proposition 4.1. *A hypergraph H is (regular) projective if and only if it has no hyperedges.*

Proof. If a hypergraph has no edges it is projective since it is a coproduct of the vertex object $I(V)$ which is clearly projective. Conversely, by the lemma above it is enough to show that each hyperedge object E_k is not projective. Let 1 be the terminal hypergraph with one vertex and one hyperedge. Then every morphism $E_k \to 1$ is a (regular) epimorphism. Let E_r be the hyperedge object with r vertices where *r* is of cardinality strictly greater than *k*. Since there is no morphism from E_k to E_r there is no factorization of $E_k \to 1$ through E_r showing E_k is not (regular) projective. \Box

For each set *X*, the interpretation functor $I: \mathbb{G}_{(X,\text{Aut}(X))} \to H$ factors through the full subcategory *H^k* of hypergraphs consisting of hypergraphs *H* such that the incidence of each edge is of cardinality less than or equal to the cardinality *k* of *X*. In other words, H_k is the slice category **Set** \downarrow P_k where *P^k* is the covariant *k*-power set functor which takes a set to the set of all subsets with cardinality less than or equal to *k*. The inclusion functor $i: H_k \hookrightarrow H$ admits a coreflector $r: H \to H_k$ which takes a hypergraph *H* to the hypergraph $r(H)$ with vertex set $r(H)(V) = H(V)$ and edge set $r(H)(E) := \{ \alpha \in H(E) \mid \# \varphi(\alpha) \leq k \}.$ Therefore the nerve realization for the interpretation $I_k :=$ $rRy: \mathbb{G}_{(X,\text{Aut}(X))} \to H_k$ is $R_k \to N_k: H_k \to \mathbb{G}_{(X,\text{Aut}(X))}$ where $R_k = rR$ and $N_k = Ni$. Moreover, by restricting to the subcategory H_k the counit $\varepsilon_H : rRN(i|H) \Rightarrow H$ is now an epimorphism.

Proposition 4.2. For a cardinal number k , the class of projective objects in H_k are precisely the *coproducts of* $I_k(V)$ and $I_k(A)$ where $I_k: \mathbb{G}_{(X, \text{Aut}(X))} \to H_k$ is the interpretation functor described *above.*

Proof. It is clear that $I_k(V)$ is projective. Since a hypergraph *H* in H_k has an edge if and only if it admits a morphism from $I_k(A)$ to it and since epimorphisms in H_k are those morphisms surjective on vertex and edge sets, it is clear $I_k(A)$ is projective. Therefore the coproducts of $I_k(V)$ and $I_k(A)$ are projective. Conversely, consider the following composition

$$
R_k(\mathrm{Proj}(N(H))) \longrightarrow R_k N_k(H) \xrightarrow{\varepsilon_H} H
$$

where the morphism $R_k(\text{Proj}(N(H))) \longrightarrow R_k N_k(H)$ is the application of R_k on the epimorphism $Proj(N(H)) \to N(H)$ described in Section 3. Note that since R_k is a left adjoint, it preserves epimorphisms. Moreover, R_k preserves colimits, therefore $R_k(\text{Proj}(N(H))) = \bigsqcup_{N(H)(V)} R_k \underline{V} \sqcup$ $\bigsqcup_{N(H)(A)} R_k \underline{A}$. Then since $R_k y = I_k$ where *y* is the Yoneda embedding, we have $R_k(\text{Proj}(N(H)))$ = $\Box_{N(H)(V)} R_k I_k(V) \sqcup \Box_{N(H)(A)} I_k(A)$. Therefore, every object in H_k admits an epimorphism from a projec[tiv](#page-4-2)e object (i.e., H_k has enough projectives). Thus every projective in H_k is a split subobject of the essential image of the functor R_k Proj: $\widehat{\mathbb{G}}_{(X,\text{Aut}(X))} \to H_k$. However it is clear that the only split subobjects are coproducts of $I_k(V)$ and $I_k(A)$.

 \Box

Proposition 4.3. *A hypergraph Q is injective if and only if Q is non-initial and for each subset of* $S \subseteq Q(V)$ *, there is an edge* $\alpha \in Q(E)$ *with incidence equal to S.*

Proof. Suppose *Q* is injective. For each subset $S \subseteq Q(V)$, let E_S be the hypergraph with one edge e with incident equal to $\varphi(e) = S$. Let $f: \Box_S I(V) \hookrightarrow E_S$ and $g: \Box_S I(V) \to Q$ be the inclusions of vertices. Since *Q* is injective there is a morphism $h: E_S \to Q$ which necessarily is a monomorphism. Hence *Q* must have an edge *q* with incidence $\varphi(q) = S$.

Conversely, let $f: G \to H$ be a monomorphism and $g: G \to Q$ be a morphism in the category of hypergraphs. Since *Q* is non-initial, there is a vertex $v \in Q(V)$. We define the morphism $h: H \to Q$ on vertices

$$
h_V(w):=\begin{cases} g_V(u) & \text{if}\,\, \exists u\in G(V), g_V(u)=w\\ v & \text{if}\,\, \forall u\in G(V), g_V(u)\neq w. \end{cases}
$$

Each edge *e* in *H* not in the image of f_A has incidence subset $S \subseteq H(V)$ which can be decomposed *S* \cong *S*₀ *⊔ S*₁ such that *S*₀ is in the image of *f_V* and *S*₁ is disjoint to the image of *f_V*. Then for such an edge e in $H(E)$, choose an edge $[e] \in Q(E)$ with incidence $g(f_V^{-1}(S_0)) \cup \{v\}$ where $g_V(f_V^{-1}(S_0))$ is the image of $f_V^{-1}(S_0)$ under g_V . Then we define

$$
h_E(e) := \begin{cases} g_E(b) & \text{if } \exists b \in G(E), g_E(b) = e \\ [e] & \text{if } \forall b \in G(E), g_E(b) \neq e \end{cases}
$$

Then h_V and h_E describe a morphism of hypergraphs $h: H \to Q$ such that $h \circ f = g$. Therefore, *Q* is injective. П

Corollary 4.4. Let Q be a hypergraph and $X := Q(V)$. Then Q is injective if and only if $N(Q)$ *is injective as a* $(X, \text{Aut}(X))$ *-graph where N is the nerve of the interpretation* $I: \mathbb{G}_{(X, \text{Aut}(X))} \to H$ *defined above.*

Proof. Let *Q* be an injective hypergraph. By Proposition 3.1, it is enough to show that for each set map *j*: $X \to N(Q)(V)$, there is an arc $\alpha \in N(Q)(A)$ such that $\partial(\alpha) = j$. The image of *j* describes a subset *S* of vertices in *Q*. Therefore by the result above, there is a hyperedge *e* with incidence equal to *S*. Let α : $I(A) \rightarrow Q$ be the arc in $N(Q)$ corresponding to the hypergraph morphism which takes the vertex *x* to $j(x)$ for each $x \in X$ and the single [hyp](#page-4-3)eredge $a \in I(A)$ to *e*. Then $\partial(\alpha) = j$ and thus $N(Q)$ is an injective $(X, \text{Aut}(X))$ -graph.

Conversely, suppose $N(Q)$ is injective and let $S \subseteq Q(V)$ a subset of vertices. Let $j: X \to N(Q)(V)$ be a set map with image equal to *S*. Then there is an arc $\alpha \in N(Q)(A)$ with incidence $\partial_Q(\alpha) = j$. Since α corresponds to the hypergraph morphism $\alpha: I(A) \to Q$, there must be an edge $e \in Q$ such that $a \in I(A)$ is mapped to *e*, i.e., *e* has incident equal to *S*. Therefore, *Q* is an injective hypergraph. \Box

4.1.1 The Category of Π*X***-Graphs**

Let *X* and *Y* be sets. We define the symmetric *X*-power of *Y*, denoted $\mathbb{I}_{X}(Y)$, as the multiple coequalizer of $(\underline{\sigma} : \Pi_X(Y) \to \Pi_X(Y))_{\sigma \in Aut(X)}$ where $\underline{\sigma}$ is the σ -shuffle of coordinates in the product. This definition extends to a functor $\mathbf{\underline{\Pi}}_X$: **Set** \rightarrow **Set**. Note that if $j: X' \rightarrow X$ is a set map, then there is a natural transformation $\underline{\Pi}_X \Rightarrow \underline{\Pi}_{X'}$ induced by the universal mapping property of the product. In particular, when $X \to X' = 1$ is the terminal map, we have $\text{id}_{\textbf{Set}} = \underline{\Pi}_1 \Rightarrow \underline{\Pi}_X$ which we denote by η : ids_{et} $\Rightarrow \underline{\Pi}_X$.¹¹

To define an interpretation functor $I: \mathbb{G}_{(X, \text{Aut}(X))} \to G_{\mathbb{I}_X}$, we let q be the unordered set $(x)_{x \in X}$ in $\mathbb{I}_{X}(X)$. Since $\mathbb{I}_{X}(\sigma)(x)_{x\in X} = (x)_{x\in X}$ for each automap $\sigma: X \to X$, the interpretation is well-defined.

Proposition 4.4. *The interpretation* $I: \mathbb{G}_{(X, \text{Aut}(X))} \to G_{\mathbb{I}_{X}}$ *is dense.*

¹¹Note that in the case $X = 2$, the category of \mathbb{I}_{X} -graphs is the category of undirected graphs in the conventional sense in which morphisms are required to map edges to edges.

Proof. Let (E, V, φ) and (K, L, ψ) be $G_{\mathbb{I}_{X}}$ -objects and $\lambda: D \Rightarrow \Delta(K, L, \psi)$ a cocone on the diagram $D: I \downarrow (E, V, \varphi) \to G_{\underline{\Pi}_X}$. Let *e* be an edge in *E* and $f: X \to V$ be the set morphism with $\Pi_X f = \varphi(e)$. Then (Fe^*, f) : $I(A) = (1, X, \top) \to (E, V, \varphi)$ is an object in $I \downarrow (E, V, \varphi)$ and thus there is a morphism $\lambda_{(Fe^i,f)} =: (Fe'^i, g): D(e^i, f) = (1, X, \top) \rightarrow (K, L, \psi)$. By the compatibility of the cocone, this gives us a uniquely defined $h: E \to K$, $e \mapsto e'$ on edges. Similarly for each vertex $v \in V$, there is a morphism $(l_E, \ulcorner v)$: $I(V) = (\emptyset, 1, \ulcorner_1) \to (E, V, \varphi)$ and a cocone inclusion $(l_K, \lceil w \rceil)$: $D(l_E, \lceil v \rceil) = (\emptyset, 1, l_1) \rightarrow (K, L, \psi)$ giving us a factorization on vertices $k: V \rightarrow L$. Since $\psi \circ h(e) = \underline{\Pi}_X(kf) \circ \top = \underline{\Pi}_X(k) \circ \varphi(e)$ for each edge *E*, (h, k) : $(E, V, \varphi) \to (K, L, \psi)$ is a well-defined $G_{\mathbf{\Pi}_X}$ -morphism which necessarily is the unique factorization of the cocone. Therefore, *I* is dense. \Box

Corollary 4.5. *The nerve* $N: G_{\mathbb{I}\mathbb{I}_X} \to \widehat{\mathbb{G}}_{(X, \text{Aut}(X))}$ *is full and faithful.*

Note that the realization functor takes a $\widehat{\mathbb{G}}_{(X,\text{Aut}(X))}$ -object and quotients out the set of arcs by Aut(X). Hence the unit of the adjunction $\eta_P: P \to NR(P)$ is bijective on vertices and surjective on arcs. Hence the adjunction is epi-reflective.

For a $G_{\underline{\Pi}_X}$ -object (B, C, φ) , the embedding given by the nerve functor is given by

$$
\begin{aligned} N(B,C,\varphi)(V) &= G_{\underline{\Pi}_X}(I(V), (B,C,\varphi)) \cong C, \\ N(B,C,\varphi)(A) &= G_{\underline{\Pi}_X}(I(A), (B,C,\varphi)) \\ &= \{\, (e,g) \, \mid \, e \in B, \ g \colon X \to C \ s.t. \ \underline{\Pi}_X g = \varphi(e) \, \} \end{aligned}
$$

The right-actions are by precomposition, i.e., $(e, g) \cdot x = (e, g \circ \tau x)$, $(e, g) \cdot \sigma = (e, g \circ \sigma)$.

Next, we show that injective and projectives in the category of $G_{\underline{\Pi}_X}$ -graphs are precisely those objects which are taken to injective and projective objects in the category of (*X,* Aut(*X*))-graphs.

Proposition 4.5. *A* \prod_X -graph *Q* is injective if and only if $N(Q)$ is an injective $(X, \text{Aut}(X))$ -graph.

Proof. If *N*(*Q*) is injective, then *Q* is injective since *N* is full and faithful and preserves monomorphisms. Conversely, let *Q* be an injective \mathbb{I}_{X} -graph and consider the monomorphism $f: G \to H$ and morphism $g: G \to N(Q)$ of $(X, \text{Aut}(X))$ -graphs. The realization preserves monomorphisms, hence $R(f): R(G) \to R(H)$ is a monomorphism. Since the counit $\varepsilon_Q: RN(Q) \to Q$ is an isomorphism, $RN(Q)$ is injective and thus there is a morphism $h: R(H) \to RN(Q)$ such that $h \circ R(f) = R(g)$. Therefore, the following diagram commutes

where $\overline{h} := \varepsilon_Q \circ N(h) \circ \eta_H$. Thus, $\overline{h} \circ f = g$ and hence $N(Q)$ is injective.

Proposition 4.6. *A* \mathbb{I}_{X} *-graph P is projective if and only if* $N(P)$ *is a projective* $(X, \text{Aut}(X))$ *graph.*

Proof. If *N*(*P*) is projective, then *P* is projective since *N* is full and faithful and preserves epimorphisms. Conversely, let *P* be a projective \mathbb{I}_{X} -graph. It is clear that $I(V)$ and $I(A)$ are projective objects in $G_{\mathbb{I}_{X}}$, thus $R(\text{Proj}(N(P))) \cong \bigsqcup_{N(P)(V)} I(V) \sqcup \bigsqcup_{N(P)(A)} I(A)$ is projective. Since the projective refinement.

 $Proj(N(P)) \rightarrow N(P)$ (see Section 3) is an epimorphism and $\varepsilon_P : RN(P) \rightarrow P$ is an isomorphism, the composition $R(\text{Proj}(N(P))) \to RN(P) \to P$ is an epimorphism. Thus, P is a split subobject of a coproduct of $I(V)$ and $I(A)$. However, the only split subobjects of such a coproduct is itself a coproduct of $I(V)$ and $I(A)$. Then since N preserves coproducts and $NI(V) = V$ and $NI(A) = A$, *N*(*P*) is projective. \Box

5 Conclusion

The categories of (X, M) -graphs we introduced in this paper are a step towards a "universal graph" theory", where questions about the various kinds of graphs can be posed in a general way. The formal structures contained in graph and hypergraph theory can be investigated in this general setting thus unifying and simplifying many results and constructions found in their separate settings. In this paper, we looked at one example, namely the construction of injective hulls and projective covers of objects. However, many other categorical constructions can be obtained by the methods introduced in this paper. For example, in Schmidt [11] and [12], we construct exponential objects in categories of (*X, M*)-graphs.

The categories of presheaves we use can be thought of as a category of variable sets which two levels (determined by the objects of the (*X, M*)-graph theories) of variability. This allows us to talk about the categories of (X, M) -graphs as (intuitionistic) s[et t](#page-12-6)heorie[s.](#page-12-7)

It can also be shown that the metacategory of (X, M) -graph theories is equivalent to the category of monoid actions on sets. Thus each monoid action on a set gives rise to a category which can be used to investigate monoid actions. This is similar to associating a presheaf category to an algebraic variety in algebraic geometry. Thus we are able to connect algebra (monoid actions) with logic (intuitionistic set theory) and geometry (graph theory) in a new way.

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Competing Interests

Author has declared that no competing interests exist.

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